

# Revisiting Practical and Usable Coercion-Resistant Remote E-Voting

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Peter Roenne, Peter Y.A. Ryan, **Najmeh Soroush**

**E-Vote-ID 2020**  
**October 2020**

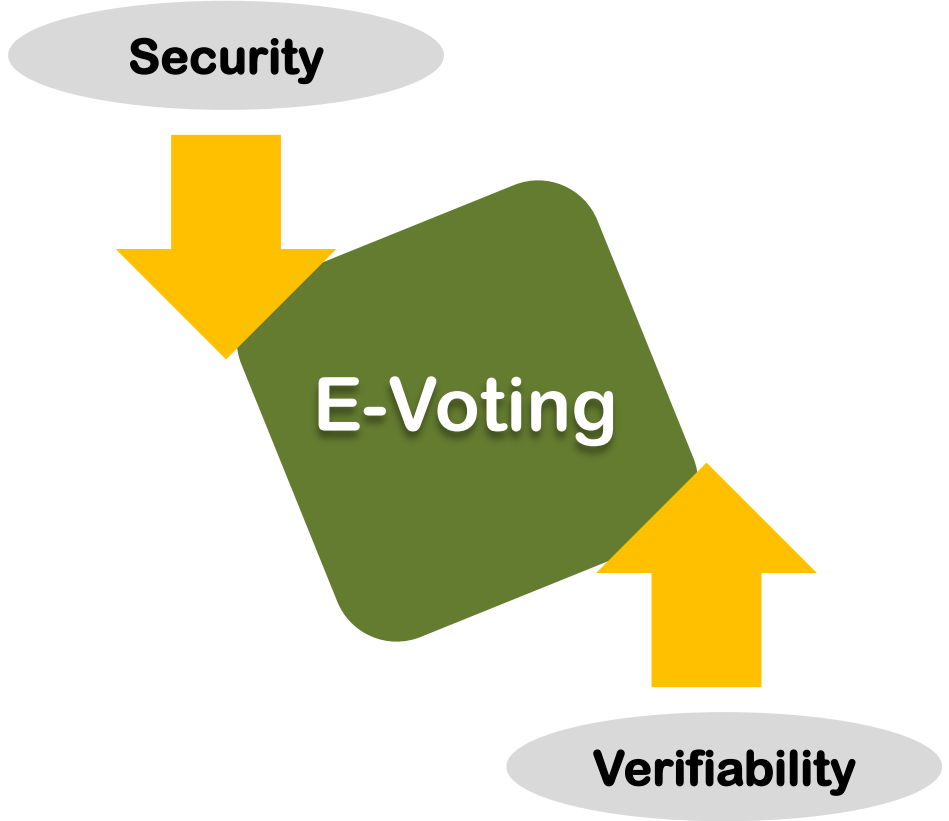


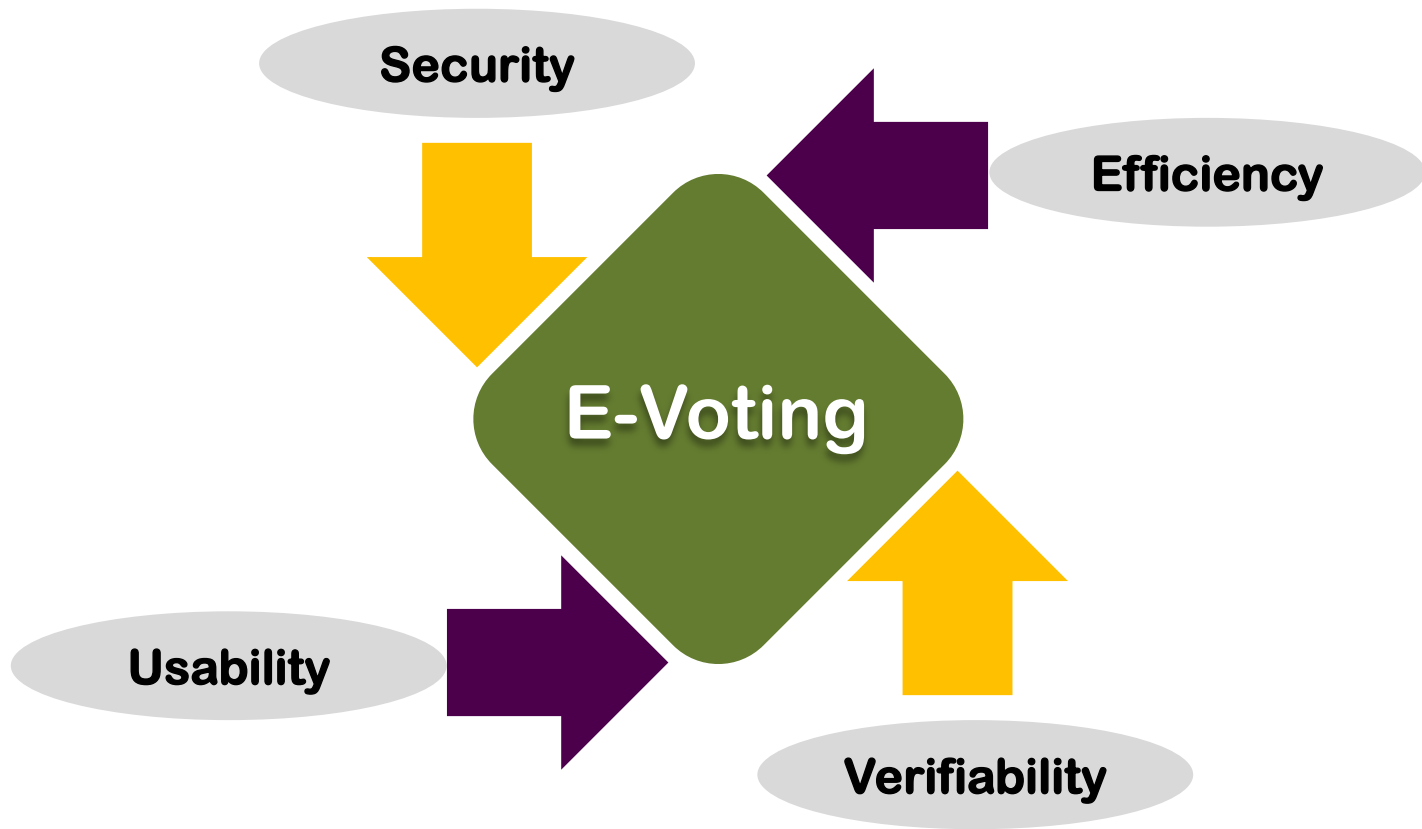
**E-Voting**

**Security**



**E-Voting**





## Participating entities:

Election Authority (EA)

Registrars

Talliers



Single

Distributed



Voter ( $V_{id}$ )

Public Bulletin Board

Secure

Verifiable

Coercion -Resistance

Receipt free

## Protocol Phases:

Set Up Phase

Key\_election=(PK,SK)

Registration Phase

Credential for legitimate voters

Voting Phase

Cast vote using credential

Tally Phase

Ballot verification

Duplicate removal

Final tally

Vérification Phase

# Usability of JCJ ?!!

Credential:  
MJ5vie9B!mj\*t3}A10PK

Long PseudoRandom string

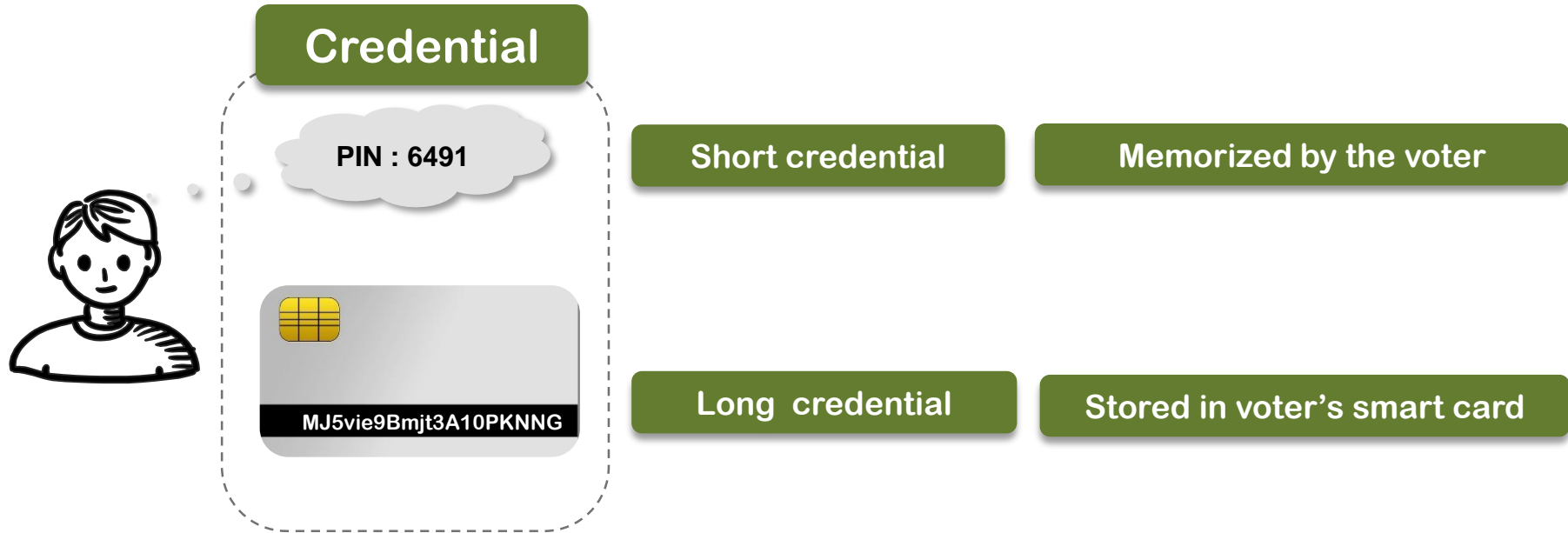
Hard to memorize by the voter,  
Storage problem

NOT human error-resilient



# Toward Usable JcJ:

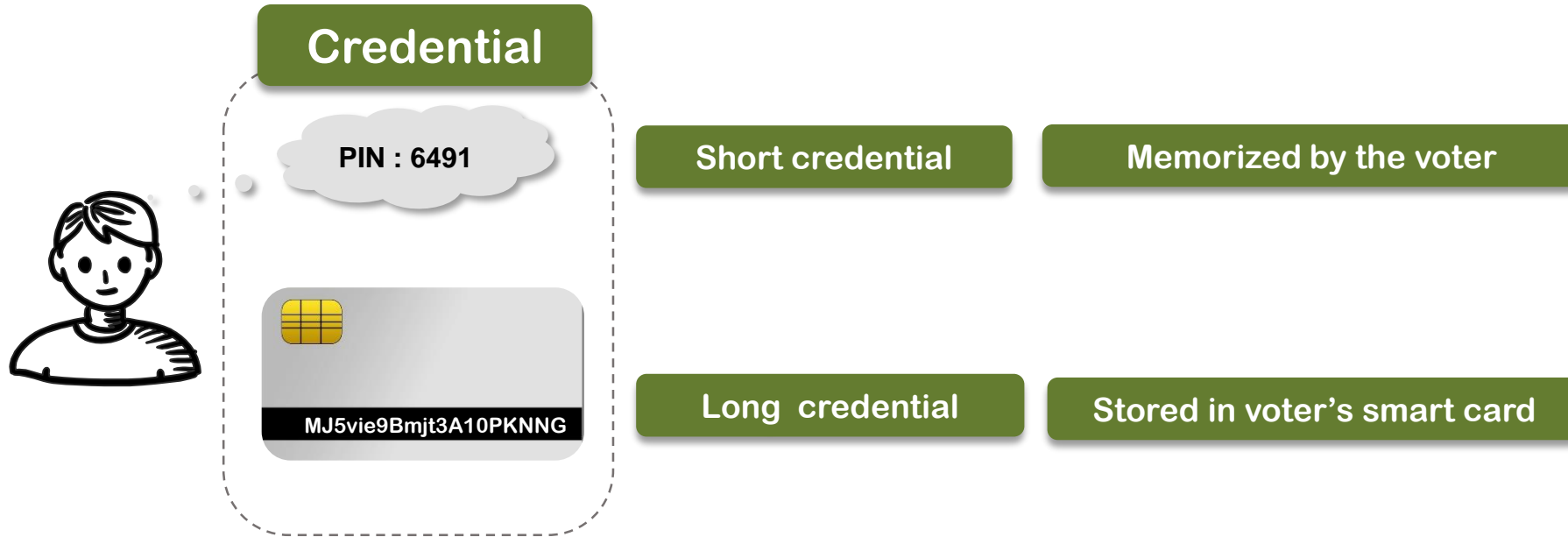
Solution by Neumann, Volkamer [NV12]:





# Toward Usable JcJ:

Solution by Neumann, Volkamer [NV12]:



Smart card removal

NOT human error-resilient

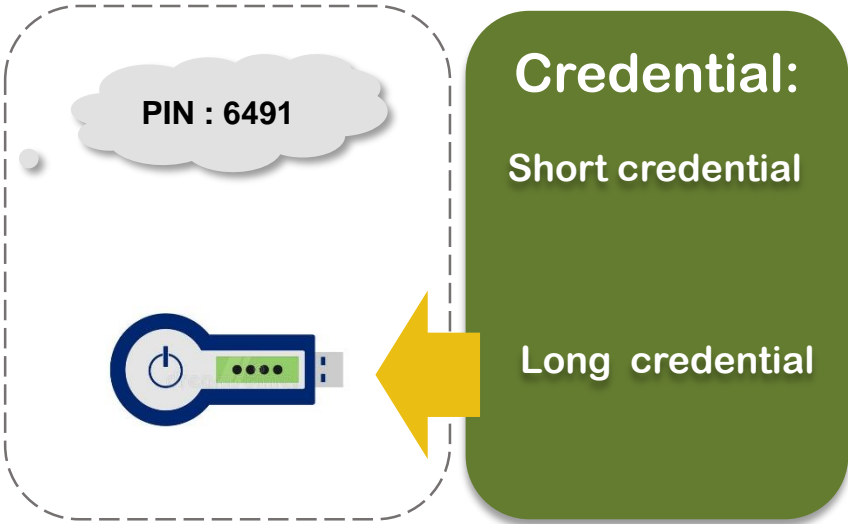
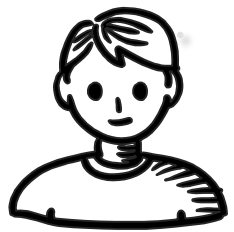
Leaky duplicate removal

Brute force attack

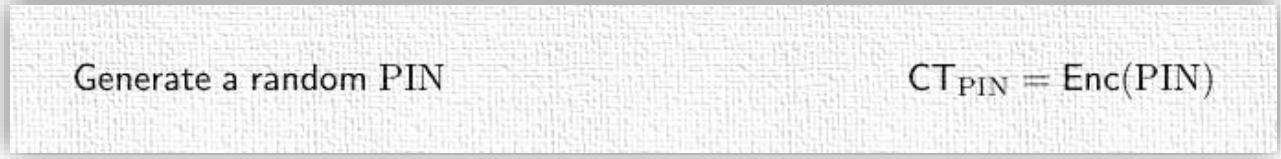
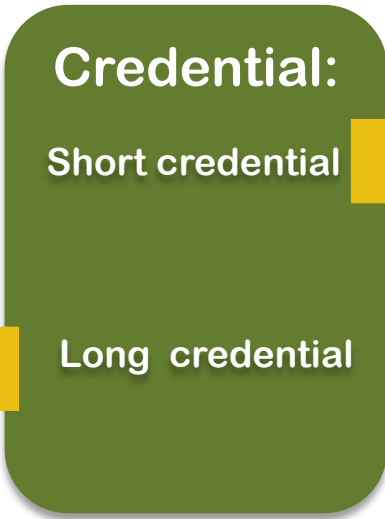
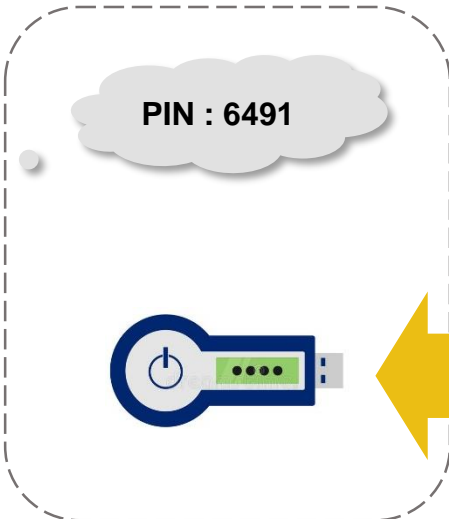
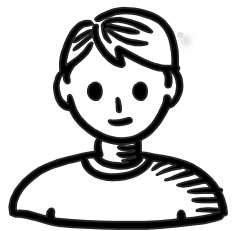
Benaloh challenge problem

Fake election identifier

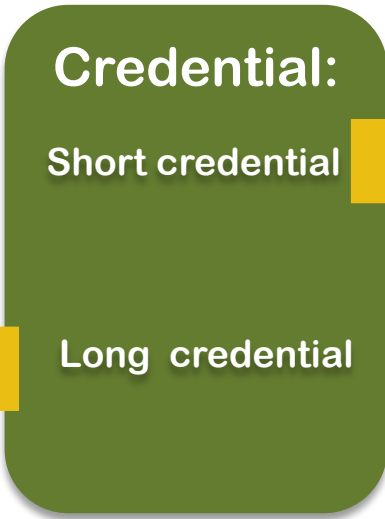
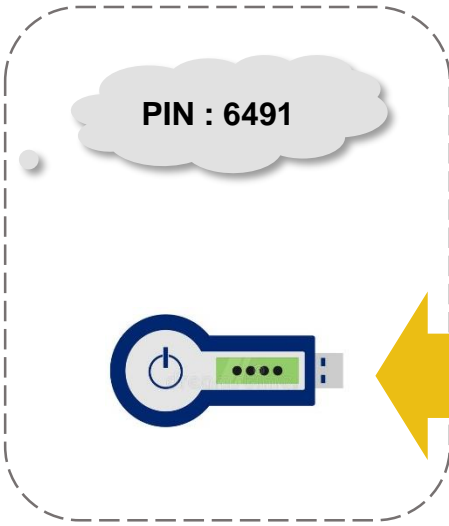
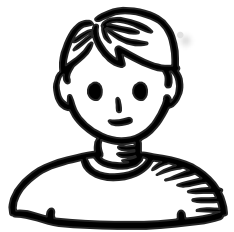
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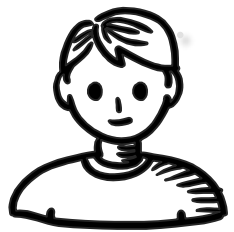
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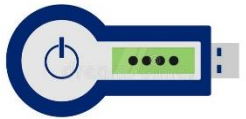
Generate a random PIN  $CT_{PIN} = Enc(PIN)$

$allowedErrorList_{\alpha} = \{\alpha_1 = \alpha, \alpha_2, \dots, \alpha_k\}$

# New solution for Usable JCJ:



PIN : 6491



Credential:

Short credential

Long credential

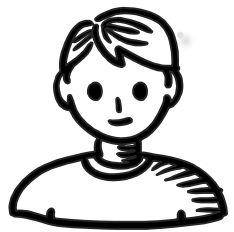
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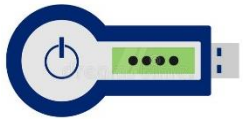
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$$\text{poly}_{PIN}(x) = \prod_{i=1}^k (x - \alpha_i) = \sum_{i=0}^k p_i x^i : (p_0, p_1, \dots, p_k)$$

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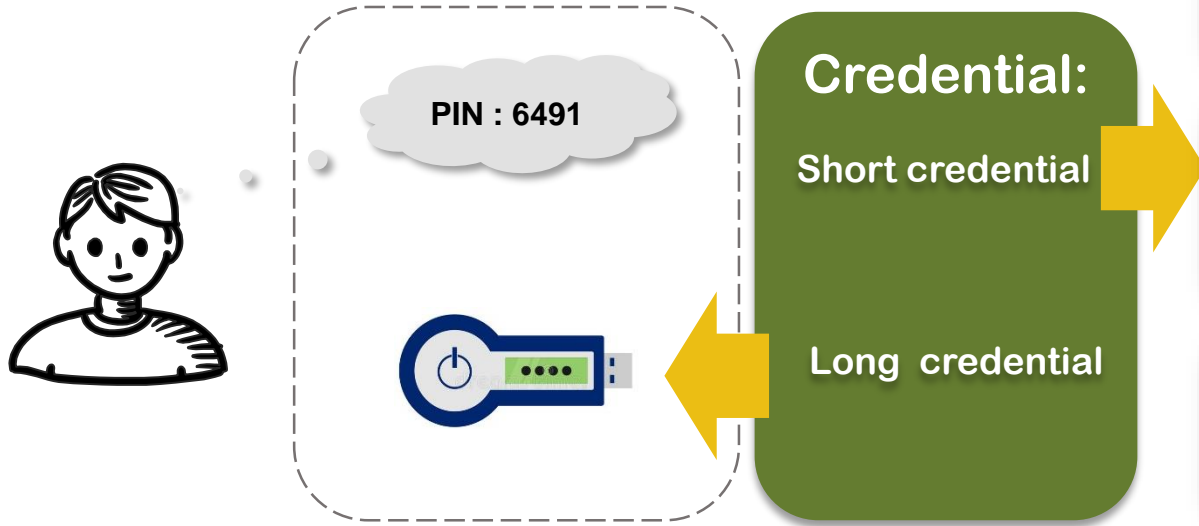
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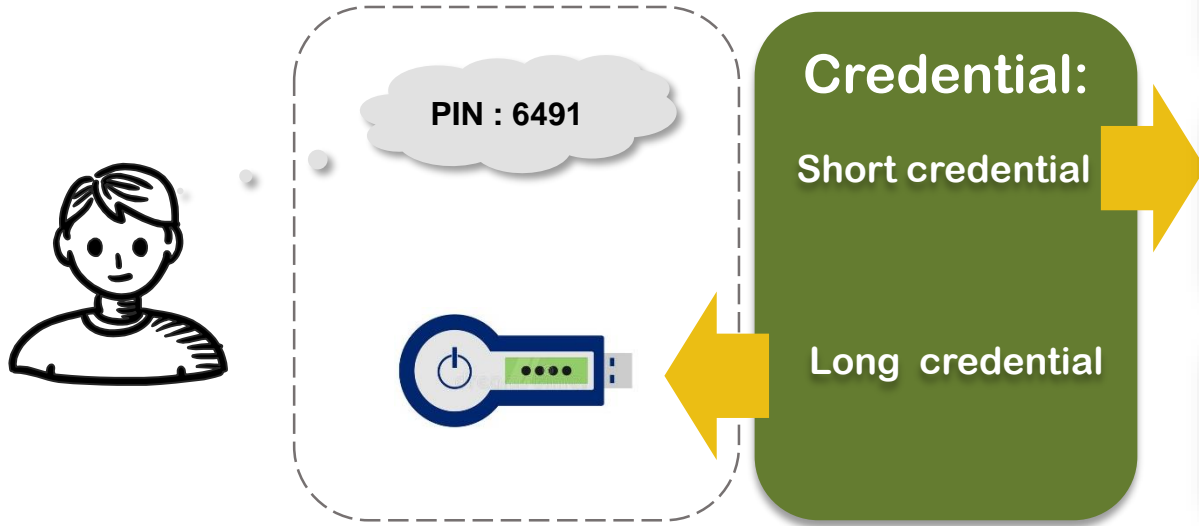
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$$\left. \begin{array}{l} CT_{PIN} = \text{Enc}(\alpha^*) \\ \text{Enc}(\text{poly}_{PIN}(x)) = \sum_{i=0}^k cp_i x^i \end{array} \right\} \stackrel{?}{\Rightarrow} (\alpha^* \in \text{ErrorList}_\alpha) \equiv \text{TRUE/FALSE}$$

$$\text{poly}_{PIN}(CT_{PIN}) = \text{poly}_{PIN}(\text{Enc}(\alpha^*)) = \text{Enc}(\text{poly}_{PIN}(\alpha^*))$$

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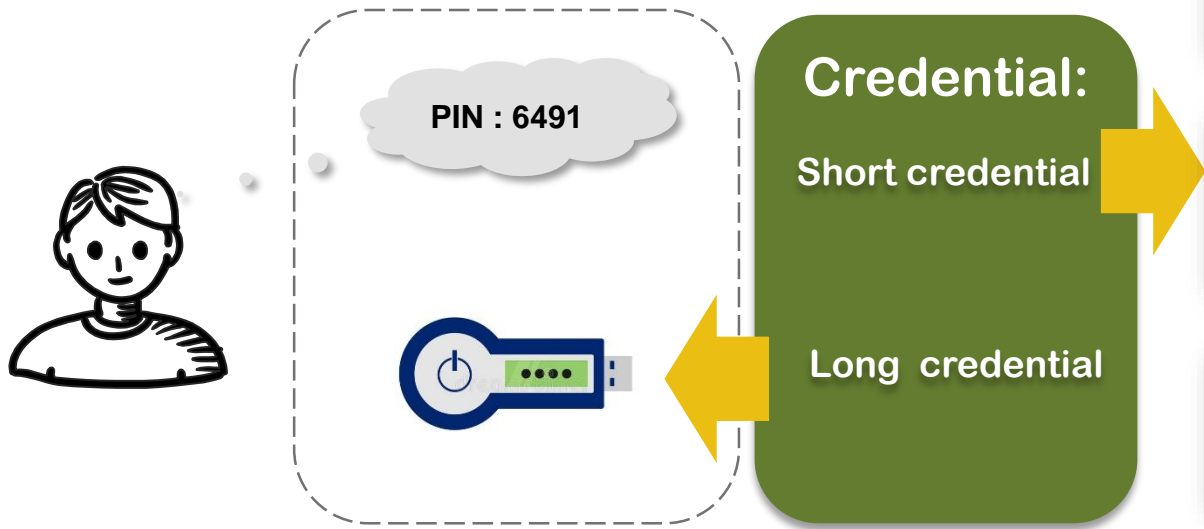


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$$poly_{PIN}(CT_{PIN}) = poly_{PIN}(Enc(\alpha^*)) = Enc(poly_{PIN}(\alpha^*))$$

- Polynomial evaluation without decryption
- Proof of the well-formedness of the polynomial
- Detect and remove ballots with invalid PIN
- Detect and remove duplicate valid ballots (valid PIN)



## Paillier Instantiation :

### Paillier Cryptosystem:

- $pk = (n = pq, \mathbb{G}, g), sk = (p, q)$
- $Enc(m) = g^m \cdot r^n \pmod{n^2}$

A partially homomorphic Encryption scheme

Security : Decisional composite residuosity assumption

Proof system: Non-Interactive sigma protocol

Evaluate the polynomial without decrypting

Efficient multi-party computation to sort ciphertext

## BGN Instantiation :

### BGN Cryptosystem:

- $pk = (n, \mathbb{G}, \mathbb{G}_T, e, g, h = g'^q), sk = (p, q)$
- $\mathbb{G} = \langle g \rangle, n = pq, e : \mathbb{G} \times \mathbb{G} \rightarrow \mathbb{G}_T$
- $Enc(m) = g^m h^r \in \mathbb{G}, m \in [T]$

A partially homomorphic Encryption scheme

Security : Discrete log and factorization

Proof system: Groth-Sahai NIWI (bilinear map)

Evaluate the polynomial without decrypting (bilinear map)

# Security Analysis concerning PIN length:

**S: Swapping errors**

PIN= 1 2 3 4 : 13 2 4, 1 2 4 3 ∈ AllowedErrorList

**W: single Wrong digit errors**

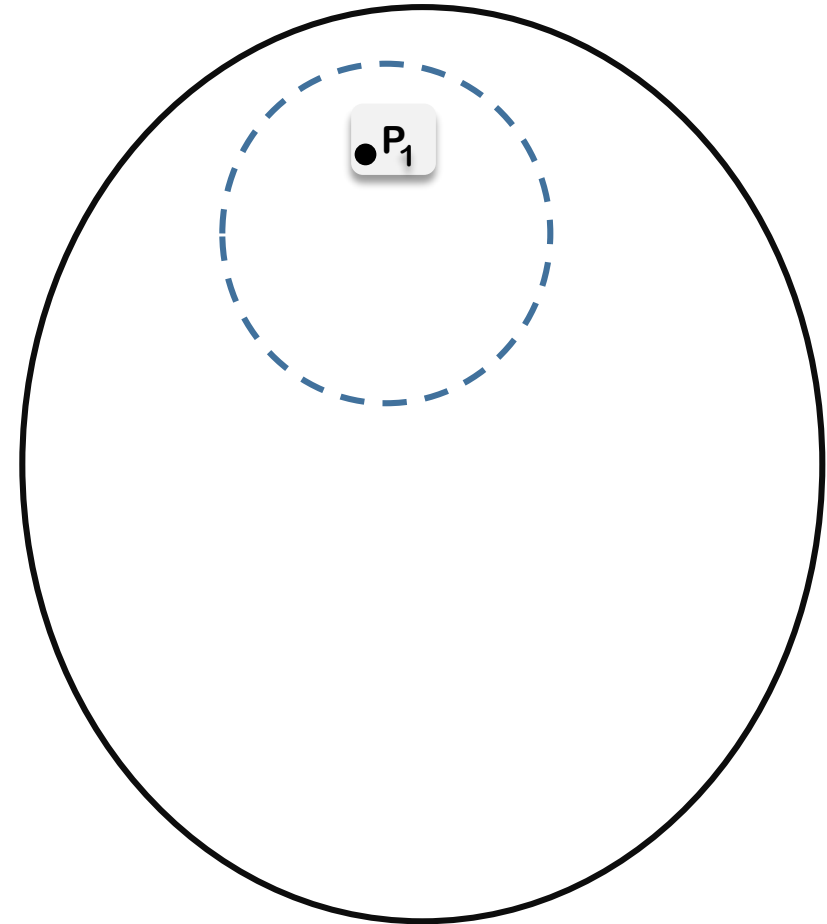
PIN= 1234 : 1235, 1434 ∈ AllowedErrorList

PIN= 1 2 3 4 :

PINs covered by “1234” : 2134, 1324 , 1243 , 1230,  
1231,1232,1233,.....,

PINs NOT covered by “1234” : 8734, 9876 , 0932 ,  
1650, 1839,1030,1891,.....,

PIN Space



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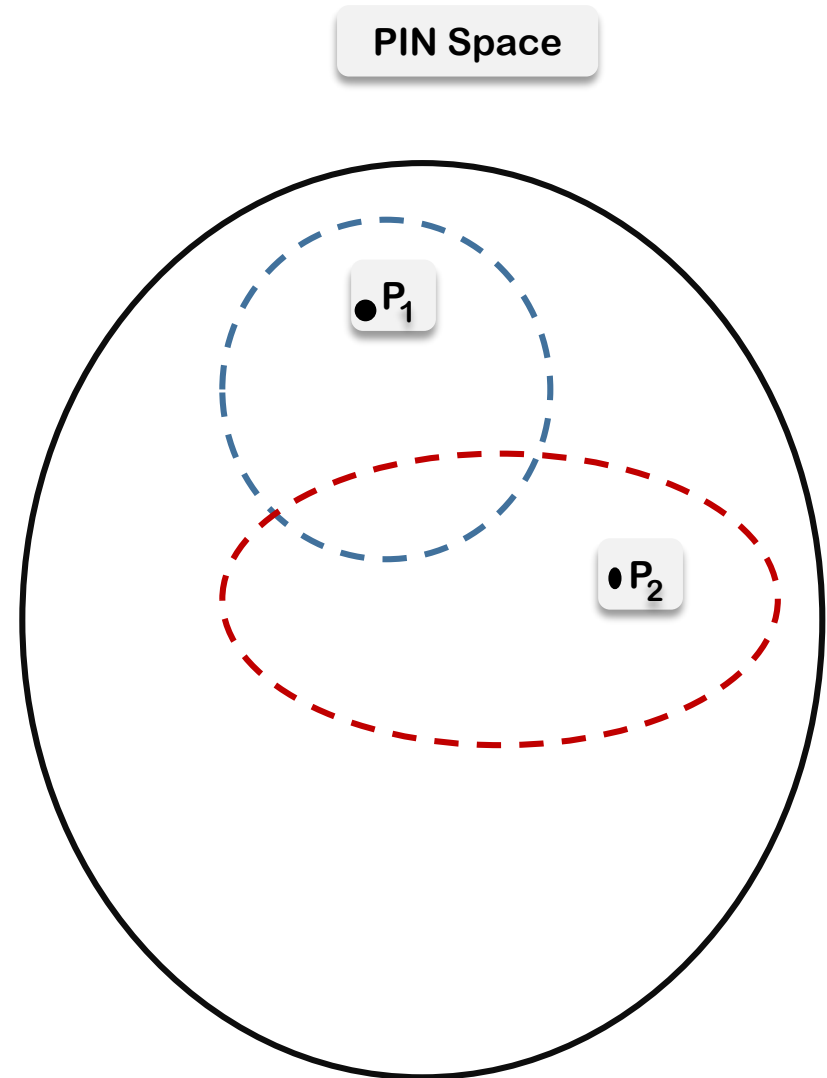
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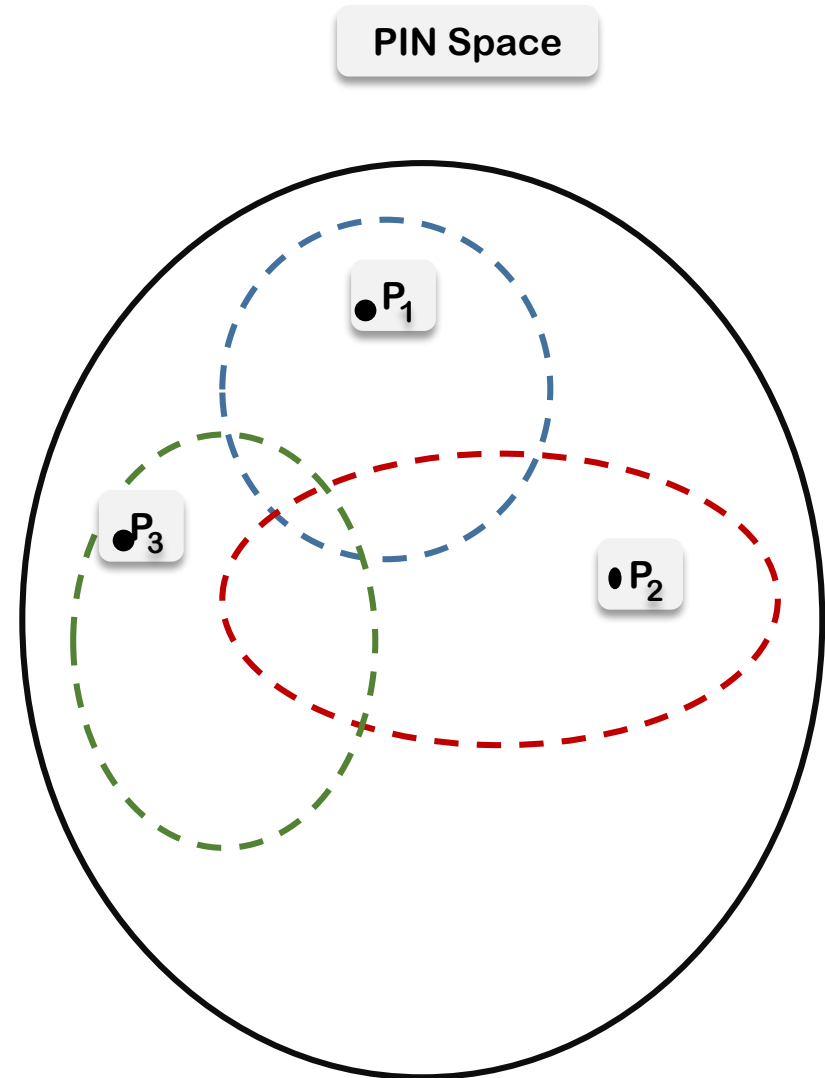
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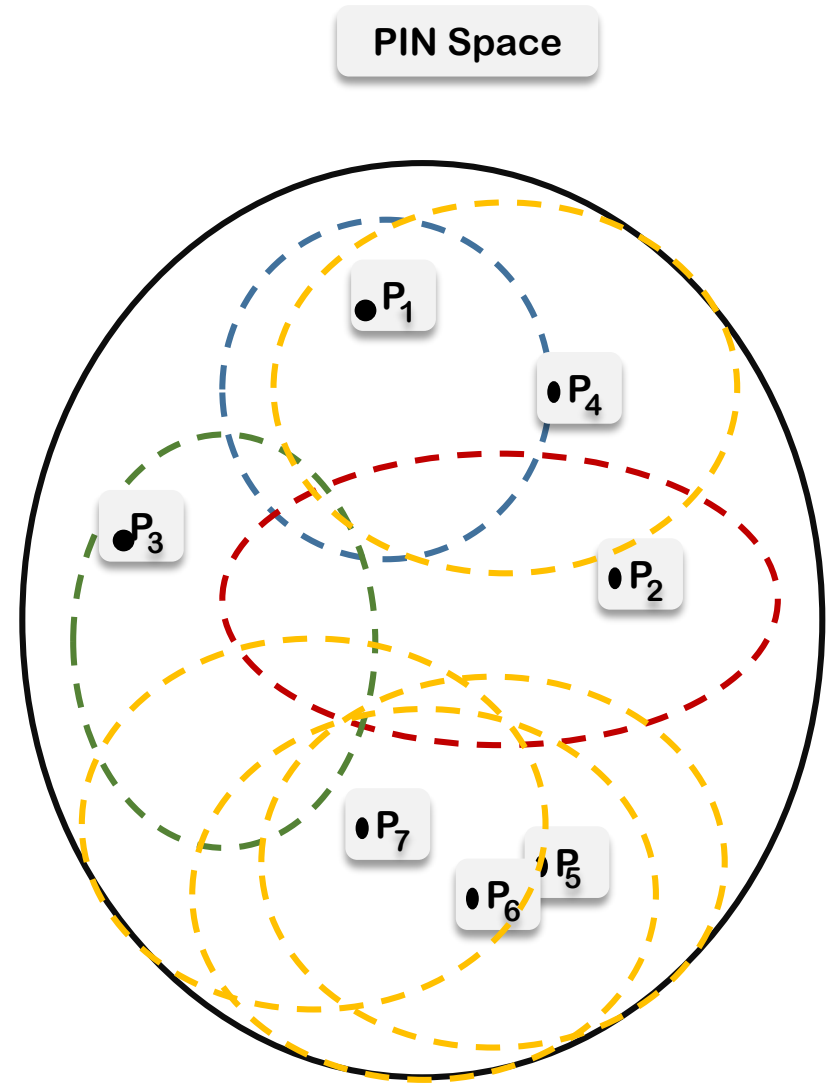
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# Security Analysis, PIN length:

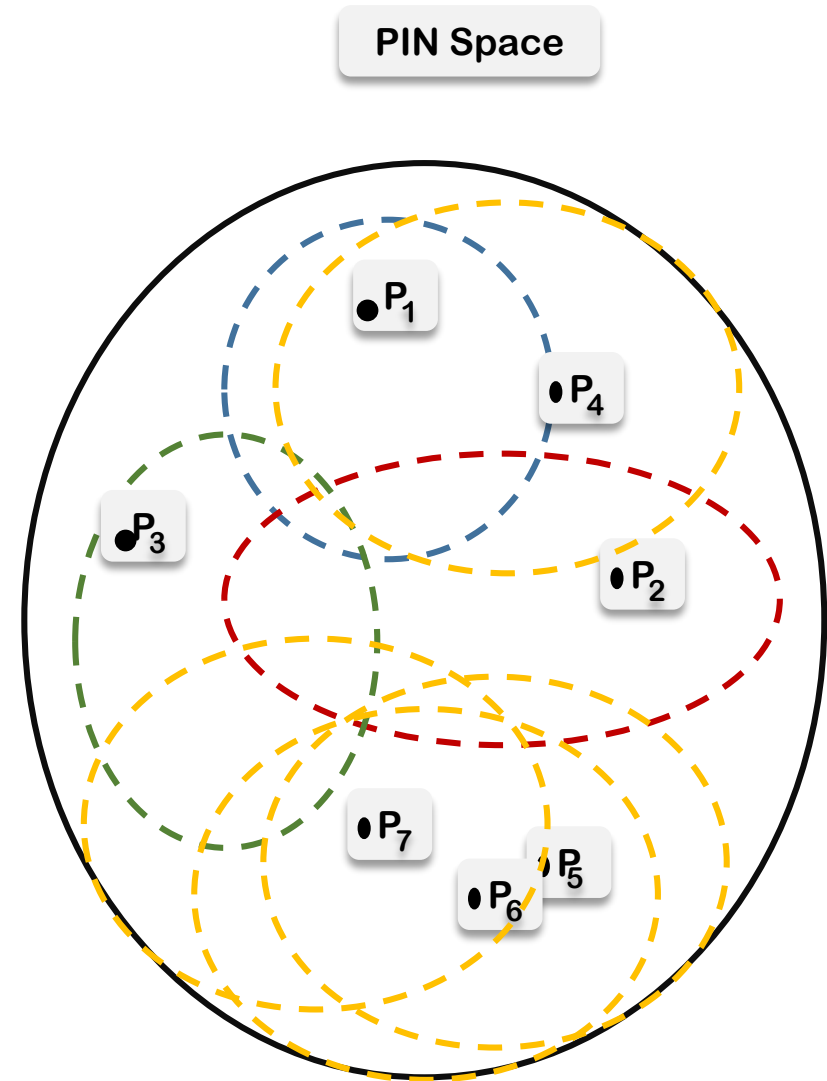
**S: Swapping errors**

PIN= 1 2 3 4 : 13 24, 1 2 43 ∈ AllowedErrorList

**W: single Wrong digit errors**

PIN= 1234 : 1235, 1434 ∈ AllowedErrorList

PIN Length	2	3	4	5
S + W Lower Bound	8	34	250	2000
S + W Upper Bound	9	78	713	6490
S Lower Bound	55	465	4131	



## Conclusions:

Presented attacks and repairs on the NV12 scheme

Presented protocols which are resilient to human errors in the form of PIN typos

## Outlook:

The digitally stored key could be combined or replaced with a key derived from biometric data

Make the error correction efficient that we can allow using noisy biometric data without fuzzy extraction.

PIN/Credential update for different elections

## Socio-technical research questions:

what is the optimal PIN policy that corrects as many PIN typos while still keeping the entropy of the PIN space sufficiently high.

Which type of PIN errors do voters do when they are in a vote setting and do not get any feedback on the correctness of the PIN.

If we do not use a smart card, or use both a smart card and key storage: how well can voters be trained to handle, fake and hide secret keys



**Thanks for your attention!**



*we're all in this together*